# *k*-block versus 1-block parallel addition in non-standard numeration systems

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## **Preliminaries**

#### Positional numeration system

Base  $\beta$  and digit set A, where

- $\beta \in \mathbb{C}$ ,  $|\beta| > 1$ , algebraic number
- $\bullet$   $\mathcal{A}\subset\mathbb{Z},$  finite set of contiguous integers containing 0

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$$z = \sum_{k=-m}^{n} a_k \beta^k$$
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$$\operatorname{Fin}_{\mathcal{A}}(eta) = \Big\{ \sum_{j \in I} x_j eta^j : I \subset \mathbb{Z}, I \text{ finite}, x_j \in \mathcal{A} \Big\}$$

## Parallel addition

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## Parallel addition – example

Impossible with  $\beta = 10$ ,  $\mathcal{A} = \{0, \dots, 9\}$ :

- $99(9)^n97 + 2 = 99(9)^n99$
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X	$\mapsto$		2	5		5		6	0	3
У	$\mapsto$	5	1	2	2	5	4	0	6	5
Z	$\mapsto$	5	3	7	4	10	9	6	6	8
0	$\mapsto$		1	<del>10</del>						
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$$p = 2$$

#### Parallel addition – known results

#### Theorem (C. Frougny, E. Pelantová, M. Svobodová, 2011)

Let  $\beta$  be an algebraic number such that  $|\beta|>1$  and all its conjugates in modulus differ from 1. Then there exists an alphabet  $\mathcal{A}\subset\mathbb{Z}$  such that addition on  $\operatorname{Fin}_{\mathcal{A}}(\beta)$  can be performed in parallel.

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#### Theorem (CF, EP, MS, 2013)

Let  $\beta$ ,  $|\beta| > 1$ , be an algebraic integer with minimal polynomial f. Let  $\mathcal A$  be an alphabet of contiguous integers containing 0 and 1. If addition in  $\operatorname{Fin}_{\mathcal A}(\beta)$  is computable in parallel, then  $\#\mathcal A \ge |f(1)|$ . If moreover  $\beta$  is a positive real number, then  $\#\mathcal A \ge |f(1)| + 2$ .

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We require  $\#A \ge |f(1)|$ ; if  $\beta$  is positive real, then  $\#A \ge |f(1)| + 2$ .

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Large increase in cardinality of alphabet may be necessary for parallelism. Potential solution: *k*-block.



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- $\bullet$  consider digits clustered in blocks of length k,

$$x = x_n \cdots \underbrace{x_{jk+k-1} \cdots x_{jk}}_{X_j} x_{jk-1} \cdots x_k \underbrace{x_{k-1} \cdots x_0}_{X_0} \bullet x_{-1} \cdots x_{-m}$$

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$$x \in \operatorname{Fin}_{\mathcal{A}}(\beta)$$
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$y \in \operatorname{Fin}_{\mathcal{A}}(\beta)$	$\ldots Y_{j+t} \ldots Y_{j+1} Y_j Y_{j-1} \ldots Y_{j-s} \ldots$	$Y_j \in (\mathcal{A})^k$
$W_j = X_j + Y_j$	$\ldots W_{j+t} \ldots W_{j+1} W_j W_{j-1} \ldots W_{j-s} \ldots$	$W_j \in (\mathcal{A} + \mathcal{A})^k$

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7 (())		
$Z_j = \phi(W_{j+t} \dots W_{j-s})$	$\ldots Z_{j+t} \ldots Z_{j+1} Z_j Z_{j-1} \ldots Z_{j-s} \ldots$	$Z_j \in (\mathcal{A})^{\kappa}$

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3 1 ( 311 3 2)	311 31 3 3 1	, , , ,

#### Compare with 1-block:

$x \in \operatorname{Fin}_{\mathcal{A}}(\beta)$	$\ldots x_{j+t} \ldots x_{j+1} x_j x_{j-1} \ldots x_{j-s} \ldots$	$x_j \in \mathcal{A}$
$y \in \operatorname{Fin}_{\mathcal{A}}(\beta)$	$\dots y_{j+t} \dots y_{j+1} y_j y_{j-1} \dots y_{j-s} \dots$	$y_j \in \mathcal{A}$
$w_j = x_j + y_j$	$\dots W_{j+t} \dots W_{j+1} W_j W_{j-1} \dots W_{j-s} \dots$	$w_j \in \mathcal{A} + \mathcal{A}$
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#### Proposition

Let  $\beta\in\mathbb{C}$ ,  $|\beta|>1$  be an algebraic integer with conjugate  $\gamma$  of modulus  $|\gamma|=1$  and let  $\mathcal{A}\subset\mathbb{Z}$  be a finite alphabet. Then no k-block p-local function can perform parallel addition on alphabet  $\mathcal{A}$ .

#### Theorem (CF, PH, EP, MS, 2013)

Given a base  $\beta$  and an alphabet  $\mathcal{B}$ . Let us suppose that there exist positive integers  $\ell$  and r such that for any  $x=x_n\dots x_0 \bullet$  and  $y=y_n\dots y_0 \bullet$  from  $\operatorname{fin}_{\mathcal{B}}(\beta)$  the sum x+y has a representation in the form

$$z = x + y = z_{n+\ell} \dots z_0 \bullet z_{-1} \dots z_{-r}.$$

Then there exists k-block 3-local function performing parallel addition in the alphabet  $\mathcal{A} = \mathcal{B} + \mathcal{B}$ , where  $k = 2(\ell + r)$ .

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Take  $\beta > 1$ , tribonacci base, i.e. root of  $x^3 = x^2 + x + 1$ .

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#### Proposition (CF, PH, EP, MS, 2013)

Let  $\beta>1$  be a number with Property (PF). Then there exists k such that k-block parallel addition is possible on the alphabet  $\{0,1,\ldots,2\lfloor\beta\rfloor\}$ .

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#### Proposition (Frougny, Solomyak, 1992)

- A number  $\beta > 1$  has Property (PF) if  $d_{\beta}(1) = \bullet t_1 t_2 \cdots t_m$  and  $t_1 \geq t_2 \geq \cdots \geq t_m \geq 1$ .
- $\beta$  has Property (PF) if  $d_{\beta}(1) = \bullet t_1 t_2 \cdots t_m t^{\omega}$  and  $t_1 \geq t_2 \geq \cdots \geq t_m \geq t \geq 1$ .



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Let  $\beta>1$  be a number with Property (PF). Then there exists k such that k-block parallel addition is possible on the alphabet  $\{0,1,\ldots,2\lfloor\beta\rfloor\}$ .

#### Proposition (Frougny, Solomyak, 1992)

- A number  $\beta > 1$  has Property (PF) if  $d_{\beta}(1) = \bullet t_1 t_2 \cdots t_m$  and  $t_1 \geq t_2 \geq \cdots \geq t_m \geq 1$ .
- $\beta$  has Property (PF) if  $d_{\beta}(1) = \bullet t_1 t_2 \cdots t_m t^{\omega}$  and  $t_1 \geq t_2 \geq \cdots \geq t_m \geq t \geq 1$ .



#### Theorem (CF, PH, EP, MS, 2013)

Let  $d_{\beta}(1) = t_1 t_2 \cdots t_m$  with  $1 \leq t_m \leq t_i$  for  $i = 2, 3, \ldots, m$  and let  $k \in \mathbb{N}$ . If parallel addition can be performed by a k-block local function on  $\mathcal{A} = \{0, 1, \ldots, M\}$ , then  $M \geq t_1 + t_m$ .

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#### Corollary

Let  $d_{\beta}(1)=t_1t_2\ldots t_m$  with  $t_1\geq t_2\geq t_2\geq \ldots \geq t_m\geq t\geq 1$  be the Rényi expansion of 1. Then there exists  $M\in\mathbb{N}$  such that parallel addition by a k-block local function is possible on the alphabet  $\{0,1,\ldots,M\}$  and  $t_1+t_m\leq M\leq 2t_1$ .

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## Block parallel addition – *d*-bonacci base

Let  $d\in\mathbb{N},\ d\geq 2.$  Choose  $\beta>1$  as the real root of  $X^d=X^{d-1}+X^{d-2}+\cdots+X+1.$ 

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- 1-block parallel addition requires  $\#A \ge |f(1)| + 2 = d + 1$
- $d_{\beta}(1) = \bullet(1)^d$  and  $\lfloor \beta \rfloor = 1$ , so k-block parallel addition is possible on the alphabet  $\mathcal{A} = \{0, 1, 2\}$ . It cannot be further reduced

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- Mutual dependence of the three parameters is yet to be investigated

Thank you for attention